To: Distribution
From: John J. Bongiovanni
Date: 03/03/82
Subject: VTOCE Buffer Management

## 1. ABSTRACI

This document describes a redesign of the physical buffer manager for Volume Table of Contents Entries (VTOCEs). Under this redesign. the I/O activity to VTOCEs is reduced considerably. at the expense of wired buffer space. Existing interfaces are preserved, although some modules outside of the VTOCE physical buffer manager are modified to improve efficiency.

The current physical buffer manager is described in general terms. followed by an overview of the design. The internals of the redesigned physical buffer manager are presented in some detail. The performance trade-offs (I/O savings versus wired memory increases) are quantified.

Send comments on this MTB by one of the following means:
By Multics Mail, on MIT or System M:
Bongiovanni. Multics
By Telephone:
HVN 261-9316 or (617)-492-9316

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Under the New Storage System (post MR 4.0). a segment resides entirely on one disk volume, and it is described by a Volume Table of Contents Entry (VTOCE). which also resides on that volume. The VTOCE, like all of Gaul, is divided into three parts: activation information. the file map. and permanent information. The activation information is a representation of the Active Segment Table Entry (ASTE). and it consists of information either needed when a segment is active or likely to be modified because a segment is active. The file map describes the disk addresses assigned to each page in the segment. The permanent information contains information which is never or rarely changed. A VTOCE is 192-words long, which is three 64-word sectors (a sector being the physical addressible unit of older disks).

The VTOCE is arranged so that the activation information and the first part of the file map are in the first sector (the first part of the file map reflects the first 96 pages of the segment). The second sector contains only the file map (pages 97 through 224). The third sector contains the remainder of the file map and the permanent information.

The physical buffer manager for VTOCEs is the module vtoc_man. vtoc_man does all I/Os as 64 -word (sector) I/Os. That is, each I/O requested by vtoc_man is for a 64 -word sector. or one-third of a VTOCE, vtoc_man contains the entry get_vtoce to read one or more sectors of a VTOCE. and the entry put vtoce to write one or more sectors of a VTOCE. The caller of each of these routines supplies a buffer and a bit mask ( 3 bits) indicating which of the three sectors in the VTOCE are to be read or written. The number of sectors read or written as a result of a call to vtoc_man can be 1, 2, or 3. depending on the setting of the bit mask. Correspondingly, vtoc_man issues 1. 2, or 3 I/Os to satisfy the request.

Typically, only part of $\operatorname{VTOCE}$ is read or written. An example is segment activation. To activate a segment, the activation information in the VTOCE is required. along with as much of the file map as there are (non-null) records. The highest non-null record (the current segment length) is part of the activation information. So activation involves reading the first part of the VTOCE to get the current length. If this is 96 or less. no other VTOCE I/O is necessary, as the requisite portion of the file map has been read. If it is larger than 96. additional VTOCE I/Os are necessary. Similarly, deactivating a segment involves writing the activation information and file map back to the VTOCE. If the current segment length is 96 or smaller, only the first sector need be written.
vtoc man uses an array of 64 -word buffers in the unpaged segment vtoc buffer_seg. There may be up to 64 of these buffers (settable by the site). Each buffer can hold one VTOCE part. and the buffers are independent. When vtoc_man is called to read 3 VTOCE parts. for example. it acquires 3 buffers and then issues 3 I/Os. There is also an optimization wherein the buffers serve as a cache. so that I/Os are not initiated for a VTOCE part which is in some buffer already, left over from some previous operation.

Note that the "unit of issue", when dealing with vtoc_man. is the VTOCE part, and not the VTOCE. In a general sense, vtocman only barely realizes that the parts belong to the same logical entity (this is not entirely true, but it is a valid description of the buffering strategy).
vtoc_man was designed when main memory was a scarce system resource. It represents a very clean design aimed at conserving main memory (buffer space) without excessive I/O overhead. The careful design of the VTOCE itself has kept I/O activity to VTOCEs relatively light (typically, $10 \%$ of all disk I/O traffic is VTOCE I/O). However, as larger systems have evolved and the economics of components have changed. the assumptions behind the original design have become less valid. Main memory is much cheaper, and disk $I / O$ is a common bottleneck on large systems. By redesigning vtoc_man, VTOCE I/O traffic can be reduced at the expense of large buffer space.

## 3. DESIGN OVERVIEW

The design is quite simple. and can be summarized concisely:
vtoc buffer_seg remains an unpaged segment. It contains a site-settable number of buffers. each of which is 192-words long (i.e., each can hold an entire VTOCE). There is no particular limit to the number of such buffers. other than the amount of abs-wirable space available. vtoc buffer_seg is described in vtoc_buffer.incl.pl1. which is attached.
vtoc_man retains all existing entry points and interfaces.
All requests to vtoc man $\$ \mathrm{get}$ vtoce result in reading the entire VTOCE (unless those parts requested are already in some buffer).

All requests to vtoc_man\$put vtoce cause only those parts indicated by the bit mask to be written. Any optimization possible is done. so that as few I/Os as possible are issued. The alternative (viz., writing an entire VTOCE)
causes performance degradation at segment deactivation time, as it would be necessary to read the VTOCE before writing it back to disk (adding an extra $1 / 0$ per deactivation).

The dctl and disk_control entries to read and write a single sector are changed to entries to read and write some (supplied) number of consecutive sectors. This requires additional information to be carried in the disk queue entry. The revised queue entry can be seen in dskdcl.incl.pll, which is attached.

## 4. INTERNALS

The following is an outline the procedural flow of vtoc_man\$get_vtoce and vtoc_man\$put_vtoce. The other entries in vtoc_man are not changed substantially. The locking strategy is identical to that in the current system and is not discussed.
vtoc_man\$get vtoce
Search for a buffer which holds some part of this VTOCE (identified by Physical Volume Table (PVT) index and VTOCE index). If found and out-of-service (I/O in progress). await the completion of the $I / O$.

If it is found and not out-of-service, check whether it contains the VTOCE parts requested by the caller. If so, return them.

Find a buffer which is not out-of-service and initiate an I/O request to read the VTOCE into that buffer.

Await completion of the $I / O$, and repeat the process (asynchronous things are going on, so it is not guaranteed to be the case that the buffer is still assigned to the requested VTOCE; if it is not, which is unlikely but possible. the process is repeated).
vtoc_man\$put_vtoce
Search for a buffer which contains some part of this VTOCE. If it is out-of-service, await completion of the I/O. If no buffer is found. find a buffer which is not out-of-service.

Copy the VTOCE parts supplied into the buffer.
Initiate the $I / O$ to write the VTOCE parts to disk. Note that
there is one case where this cannot be done in a single I/O (viz. parts 1 and 3 ). This will be done by initiating one I/O. awaiting its completion. and initiating the second $1 / 0$. There is currently no supervisor module which writes parts 1 and 3.

## 5. QTHER CHANGES

The following are other changes to vtoc man. to be made at the same time, but not related to the primary objectives of the design. These changes will improve the processor efficiency of vtoc_man at a modest increase in memory.

The linear list of buffer descriptors is replaced by a doubly-threaded used list. As buffers are used, they are moved to the tail of this list. Finding a buffer involves following the used list until one is found which is not out-of-service. This replaces a linear search of buffer descriptors.

A hash table is used to determine whether a given VTOCE (as identified by a PVT index and a VTOCE index) has a buffer assigned. This replaces a linear search.

## 6. PEREORMANCE

In the words of the cashier at Colleen's Chinese Cuisine. "There's no such thing as a free lunch". And that's the case here.

Metering during peak activity on System M and MIT indicate that between $20 \%$ and $25 \%$ of all VTOCE I/Os can be eliminated with this design. This corresponds to between $1 \%$ and $2 \%$ of all disk I/Os.

The cost in wired memory is approximately 128 additional words per buffer. With the current default number of buffers (32). this is an additional 4 KW of wired memory. With the current maximum number of buffers (64). this is an additional 8 KW of wired memory.

## 7. SUMMARY OF CHANGES

All changes required to implement this design are indicated below by module.
activate
Call vtoc_man\$get_vtoce for the entire VTOCE. instead of reading part 1 to determine how many parts to read.
detl
Rework for new queue format.
device_meters
Recompile with new include files.
disk_control
Rework for new queue format.
disk_init
Recompile with new include files.
disk_meters
Recompile with new include files.
disk_queue
Rework for new queue format.
get_io_segs
Recompile with new include files.
ioi_assign_disk channels
Recompile with new include files.
hc_dmpr_primitives
Call vtoc man\$get vtoce for the entire VTOCE. instead of reading part 1 to determine how many parts to read.
init vtoc_man
Change to initialize the new vtoc buffer seg.
spg_fs_info
Recompile with new include files.
truncate vtoce
Call vtoc_man\$get vtoce for the entire VTOCE. instead of reading part 1 to determine how many parts to read.
update vtoce
Call vtoc man\$get vtoce for the entire VTOCE, instead of reading part 1 to determine how many parts to read.

MTB-570
verify_lock
Recompile for new include files.
vtoc buffer_meters
Rewrite to print new metering data.
vtoc_interrupt
Rework for new sector I/O scheme.
vtoc_man
Rewrite.
wired_shutdown
Recompile with new include files.
"BEGIN INCLUDE FILE dskdci.incl.alm
"Created 02/04/82 1712.6 est Thu by convert_inciude_file,
Version of 12/01/81 1540.3 est Tue.
"Made from >user dir_dir>Multics>Bonglovanni>htd>no_salvage_dir>dskdci. Inci.pl1, modified $02 / \overline{0} 4 / 8 \overline{2} \quad 1712.5$ est Thu

Structure disk_data
equ disk_data_size, 72
equ disk_data.subsystems.0
equ disk_data.free_offset.1 n UPPER

| equ | disk_data.status_mask,2 |
| :--- | :--- |
| equ | disk_data.array, 8 |$\quad$ " LEVEL 2

equ disk_data.offset. 8 "UPPER
equ disk_data.name. 9

Structure disktab

| equ | disktab. lock, 0 |  |
| :---: | :---: | :---: |
| equ | disktab.nchan, 1 |  |
| equ | disktab.ndrives, 2 |  |
| equ | disktab.channels_online. 3 |  |
| equ | disktab.dev_busy, 4 | " DOUBLE |
| equ | disktab.dev_queued, 6 | " DOUBLE |
| equ | disktab.wq, 8 | " LEVEL |
| equ | disktab.free_q, 10 | * LEVEL |
| equ | disktab.abs_mem_addr, 11 |  |
| equ | disktab.errors. 13 |  |
| equ | disktab.ferrors, 14 |  |
| equ | disktab.edac_errors. 15 |  |
| equ | disktab.pg_to_count. 16 |  |
| equ | disktab.vt_io_count, 18 |  |
| equ | disktab.call_lock_meters, 20 | " LeVEL |
| equ | disktab.int_lock_meters, 24 | " LEVEL |
| equ | disktab.alloc_wait_meters. 28 | " Level 2 |


| equ | disktab.run_lock_meters, 32 | " LEVEL 2 |
| :---: | :---: | :---: |
| equ | disktab.pg_wait, 36 | " DOUBLE |
| equ | disktab.vt_wait,40 | " DOUBLE |
| equ | disktab.pg_10,44 | " DOUBLE |
| equ | disktab.vt_10,48 | " DOUBLE |
| equ | disktab.quēue, 52 | " LEVEL 2 |
| equ | disktab.chantab, 308 | " Level 2 |
| equ | disktab.devtab,500 | * Level 2 |

Structure quentry

| equ | quentry_size,4 |  |  |
| :---: | :---: | :---: | :---: |
| equ | quentry.next, 0 | * | UPPER |
| equ | quentry.write_sw_word, 0 |  |  |
| bool | quentry.write_sw,400000 | " | DL |
| equ | quentry.sect_sw_word,0 |  |  |
| bool | quentry.sect_sw, 200000 | " | DL |
| equ | quentry.testing_word, 0 |  |  |
| bool | quentry.testing. 100000 | " | DL |
| equ | quentry.retry_word, 0 |  |  |
| bool | quentry.retry,040000 | * | DL |
| equ | quentry.used_word,0 |  |  |
| bool | quentry.used. 020000 | " | DL |
| equ | quentry.swap_word, 0 |  |  |
| bool | quentry.swap.010000 | " | DL |
| equ | quentry.cylinder_word.0 |  |  |
| equ | quentry.cylinder_shift,0 |  |  |
| bool | quentry.cylinder_mask,007777 |  |  |
| equ | quentry.pdi_word, 1 |  |  |
| equ | quentry.pdi_shift, 30 |  |  |
| bool | quentry.pdi_mask,000077 |  |  |
| equ | quentry.coreadd_word, 1 |  |  |
| equ | quentry.coreadd_shift,6 |  |  |
| equ | quentry.dev_word. 1 |  |  |
| equ | quentry.dev_shift,0 |  |  |
| bool | quentry.dev_mask,000077 |  |  |
| equ | quentry.sector_word, 2 |  |  |
| equ | quentry.sector_shift. 15 |  |  |
| equ | quentry.n_sectors_word, 2 |  |  |
| equ | quentry.n_sectors_shift. 9 |  |  |
| bool | quentry.n_sectors_mask,000077 |  |  |
| equ | quentry.time. 3 |  |  |
| Structure | chantab |  |  |
| equ | chantab_size. 24. |  |  |
| equ | chantab.chx. 0 |  |  |
| equ | chantab.101_ctx.1 |  |  |



## Structure devtab

equ devtab_size, 8

| equ | devtab.pvtx_word,0 |
| :--- | :--- |
| equ | devtab.pvtx_shtft, 27 |
| bool | devtab.pvtx_mask,000777 |

equ devtab.inop word. 0
bool devtab.inop,000400
equ devtab.was broken_word,o
1 devtab.was broken-word.0
equ devtab.broken word,o
devtab.broken_word,o
qu devtab.broken,000100
$\begin{array}{ll}\text { equ devtab.abandoned_word, } \\ \text { bool } & \text { devtab.abandoned, } 000040\end{array}$
equ devtab.buddy_word.o
equ devtab.buddy_shift, 6
bool devtab.buddy_mask,000077
equ devtab.pdi_word,o
equ devtab.pdi_word,o 0
bool devtab.pdi_mask,000077
qu devtab. queue count, 1
qu devtab.cylinder,2
equ devtab.seek_distance, 3
equ devtab.read_count. 4
equ devtab.write_count.5
equ devtab.time_inop, 6
" DU
" DU
" ou
" DU
" double

Structure pvtdi

| equ | pvtdi.sx_shift,24 |
| :--- | :--- |
| bool | pvtdi.sx_mask,007777 |
| equ | pvtdi.usable_sect_per_cyl_shift. 12 |
| bool | pvtdi.usable_sect_per_cyl_mask,007777 |
| equ | pvtdi.unused_sect_per_cyl_shift,o |
| bool | pvtdi.unused_sect_per_cyl_mask,007777 |

Structure disk_lock_meters
equ disk_lock_meters_size,4
equ disk_lock_meters.count,0
equ disk_lock_meters.waits. 1
equ disk_lock_meters.wait_time. 2 m DOUBLE

| equ | RST_LISTX, 1 | " MANIFEST |
| :--- | :--- | :--- |
| equ | SC LISTX,2 | " MANIFEST |
| equ | DSC LISTX, 6 | " MANIFEST |
| equ | RSSLISTX, 8 | " MANIFEST |

"END INCLUDE FILE dskdci.incl.alm

```
/* Begin include file ....... dskdcl. Incl.pll
/* Structures used by the Disk DIM */
/* format: style4,delnl,insnl,tree,ifthenstmt,indnoniterend
dc1 disk_seg$ ext;
```

dcl disksp ptr.
diskp ptr:
dc 1
disk data based (disksp) aligned,
2 subsystems fixed bin,
2 free_offset bit (18)
2 status mask bit (36),
2 pad (5) fixed bin.
2 array (32).
1
3 offset bit (18),
3 pad bit (18)
3 name char (4)
) unal;
dcl 1 disktab based (diskp) aligned,
2 lock bit (36) unal,
2 nchan fixed bin,
2 ndrives fixed bin,
2 channels_online fixed bin,
2 dev_busy bit (64),
2 dev_queued bit (64)
2 wq (0:1) like qht
free_q like qht.
2 abs_mem_addr fixed bin (26) unsigned.
2 pad-fixed bin,
2 errors fixed bin,
ferrors fixed bin.
edac errors fixed bin,
$2 \mathrm{pg} i o \mathrm{count}(0: 1)$ fixed bin,
2 vt io_count ( $0: 1$ ) fixed bin,
2 call_lock_meters like disk_lock_meters
2 int_lock_meters like disk_lock_meters.
2 alloc_wait_meters like disk_lock_meters
2 run_löck mēters like disk_lōck_méters,
2 pg _wait (0:1) fixed bin ( $\overline{5} 2$ ),
2 gg wait ( $0: 1$ ) fixed bin (52).
2 pg_io (0:1) fixed bin (52).
2 vt io (0:1) fixed bin (52).
2 queue (64) like quentry.
2 chantab (8) like chantab.
2 devtab (O refer (disktab.ndrives)) like devtab;
\%page;
dcl qp ptr.
cp ptr;
/* disk data segment */
/* pointer to disk subsystem info */
/* pointer to disk DIM info structure */
/* disk subsystem information */
/* number of subsystems */
/* number of subsystems */
/* offset of first unused location in segment */
/* mask for checking for disk error */
/*. I ine up on 0 mod 8 boundary */
/* per subsystem info */
/* location of data for this subsystem */
/* name of subsystem */
/* control structure for DIM's */
/* data base lock */
/* number of disk channels */
/* highest disk drive number */
/* number of disk channels actually in use */
/* busy bit for each device */
/* requests queued bit for each device */
/* wait queue head/tail */
/* free queue head/tail */
/* absolute memory address of this structure */
/* error count */
/* fatal error count */
/* count of EDAC correctable errors */
/* count of page $1 / 0$ operations */
/* count of VTOCE $1 / 0$ operations */
/* lock meters for call side of DIM */
/* lock meters for interrupt side of DIM */
/* meters for queue entry allocations */
/* lock meters for run calls */
/* total time spent waiting for page $1 / 0$ */
/* total time spent waiting for VTOCE $1 / 0$ */
/* total time spent doing page I/0 */
/* total time spent doing page I/O */
/* total time spent doing VTOCE I/O */
/* queue entries */
/* channel information table */
/* device information table */
/* pointer to queue entry */
/* pointer to channel information table*/
do'l 1 quentry based (qp) aligned, (

2 next bit (18)
2 write_sw bit (1)
2 sect_sw bit (1),
2 testing bit (1),
2 retry bit (1),
2 used bit (1),
2 swap bit (1),
2 cylinder fixed bin ( 11 ).
2 pdi unsigned $f$ ixed bin (6),
2 coreadd bit (24),
2 dev unsigned $f$ ixed bin ( 6 ),
2 sector bit (21),
2 n_sectors fixed bin (6) unsigned.
2 pad bit (9).
2 time bit (36)
) unal;
dcl 1 chantab based (cp) aligned.
2 chx fixed bin (35).
2 ioi ctx fixed bin (35),
2 statusp ptr.
2 chanid char (8),
$($
2 padO bit (18),
2 in_use bit (1)
2 active bit (1)
2 rsr bit (1)
2 prior bit (i).
ioi use bit (i).
inop bit (1),
2 broken bit (1)
2 action_code bit (2)
2 padi bit (9)
) unal,
2 arp bit (18).
2 pad2 bit (3),
2 command bit (6),
2 erct fixed bin (8)
) unal.
2 select_data,
(
3 limit bit (12)
3 mbz bit (3).
3 sector bit (21)
) unaligned,
2 connect_time fixed bin (52).
2 connects fixed bin,
2 detailed status ( $0: 8$ ) bit (8) unal.
2 rstdew bīt (36).
2 scdew bit (36).
2 sddew bit (36).
2 dcdew bit (36),
2 dddew bit (36)
2 dscdew bit (36).
2 dsddcw bit (36).
2 rssdcw bit (36),
2 status bit (36) aligned;
/* queue entry */
/* index to next queue entry */
/* index to next queue entry */ $/$ non-zero for write operation */
/* non-zero for single sector operation */
/* non-zero if quentry is for disk ready test */
/* non-zero if retry has been performed on broken device */
/* non-zero if queue entry in use */

```
/* disk cylinder number *
/* pdi of device */
* memory address for data transfer *
** disk device code */
/* disk device code */
```

/* disk sector address */
/* number of sectors for sector $1 / 0$ */
/* low-order microsecond clock at queue */
/* time entry was queued */
/* channel information table */
/* to_manager channel index */
/* toí channel table index */

* pointer to hardware status word */
/* channel name */

```
/* non-zero if channel being used */
/* non-zero if channel active */
/* non-zero if RSR in progress */
/* priority of current request */
/* non-zero if channel usurped by IOI */
/* non-zero if channel inoperative */
/* non-zero if channel broken */
/* saved from status */
```

/* rel ptr to queue entry */
/* peripheral command */
/* error retry count */
/* data passed to IOM on select */
/* limit on number of sectors */
/* sector address */
/* time of last connect */
/* count of connects performed */
/* detailed status bytes */
/* restore command */
/* select command */
/* select data xfer */
/* command to read or write */
/* data xfer DCW */
/* data xfer DCW */
/* RSR command */
/* RSR data xfer */
/* RSR data xfer */
/* RSS command *//
/* RSS command */
\%page;
dcl. 1 qht aligned based, 2 (head, tail) bit (18) unal;
dcl dp ptr
putaip ptr:
dcl 1 devtab based (dp) aligned,
1
2 pvtx fixed bin (8).
2 inop bit (1),
2 was_broken bit (1),
2 broken bit (1),
2 abandoned bit (i).
2 pad bit (11).
2 buddy unsigned fixed bin (6),
2 pdi unsigned fixed bin (6)
) unal.
2 queue_count fixed bin (8).
2 cylinder fixed bin (11),
2 seek distance fixed bin (35, 18)
2 read_count fixed bin.
2 write count fixed bin,
2 time inop fixed bin (52);
dcl 1 putdi based (putdip) aligned.
(
2 sx fixed bin (1i),
2 usable_sect_per_cyl fixed bin (11),
2 unused_sect_per_cyl fixed bin (11)
) unal;
dcl 1 disk_lock_meters based aligned.
2 count fixed bin,
2 waits fixed bin
2 wait_time fixed bin (52);
del 1
RST_LISTX init (1),
SC_LISTX intt (2),
DSC_LISTX init (6).
RSS_LISTX init (8)
) fixed bin (12) static options (constant)
/* End of include file ...... dskdci.incl.pli */
/* queue head/tail structure */
/* pointer to device information table */
/* pointer to dim_info in PVT entry */
/* device information table */
/* index of PVT entry for device */
/* device inoperative */
/* device previously broken */
/* device down */
/* device lost and gone forever */
/* other device on this spindle or 0 */
/* primary device index */
/* count of requests queued for device */
/* current cylinder position */

* average seek distance */
/* count of reads */
/* count of writes */
/* time drive became inoperative */
/* disk DIM info in PVT entry */
/* structure index */
/* \# of usable sectors on disk cylinder */
/* \# of unused sectors at end of cylinder */
/* lock meters for disk DIM */
* total number of attempts */
/* number of attempts which required wafting */
$/ *$ total time spent waiting */
/* listx for restore */
/* listx for select */
/* listx for RSR */
* listx for RSS */
"BEGIN INCLUDE FILE vtoc_buffer.inci.alm
"Created 02/05/82 2036. B est Fri by convert_include_file,

Version of 12/01/81 1540.3 est Tue.
"Made from >udd>Multics>Bongiovanni>hardcore_test_dir>no_salvage_dir>vtoc_buffer.incl.pli, modified 02/05/82 2036.8 est Fri

Structure vtoc buffer

| r.lock, 0 - LEVEL 2 |  |
| :---: | :---: |
| equ | $v$ toc_buffer.processid.o |
| equ | vtoc_buffer.wait_event, 1 |
| equ | vtoc_buffer.notify_sw_word, 2 |
| bool | vtoc_buffer.notify_sw, 400000 " DU |
| equ | vtoc_buffer.n_bufs, 3 |
| equ | $v$ toc_buffer.n_hash_buckets, 4 |
| equ | $v$ toc_buffer.hash_mask, 5 |
| equ | $v$ toc_buffer.abs_addr. 6 |
| equ | vtoc_buffer.wait_event_constant. 8 " DOUBLE |
| equ | vtoc_buffer.buf_desc_offset, 10 " UPPER |
| equ | vtoc_buffer.buf_offset, 11 " UPPER |
| equ | vtoc_buffer.hash_table_offset, 12 " UPPER |
| equ | vtoc_buffer.meters, 13 " Level 2 |
| equ | vtoc_buffer.hash_table, 14 " UPPER |
| equ | vtoc_buffer.buf_desc.0 " LEVEL 2 |
| equ | vtoc_buffer.buffer.0 * LEVEL 2 |
| Structure | vtoc_buf_desc |
| equ | $v$ toc_buf_desc_size. 4 |
| equ equ | vtoc_buf_desc.pvte_rel.0 " UPPER <br> vtoc_buf_desc.vtocx. 0 " LOWER |
|  |  |


/* START DF: vtoc_buffer.incl.pl1 Feruary 1982

| dc1 | vtoc_buffer_seg $\$$ | ext; |
| :--- | :--- | :--- |
| dc1 | vtoc_buffer_segp | ptr; |
| dc1 | $v t o c$ buf_descp | ptr; |
| dcl | $v t o c \_b u f p$ | ptr; |

aligned based (vtoc_buffer_segp),

2 lock.
3 processia
3 wait_event
3 notify_sw
2 n_bufs
2 n_hash_buckets
2 hash_mask
2 abs_addr

2 buf offset
2 hash table offset bit (18).
2 meter̃s.
3 pad
2 hash_table

2 buf_desc
2 buffer

1 vtoc_buf_desc
2 pvte_rel
2 vtocx
2 part_desc_rel
2 part $\bar{s}$ _used
2 parts_os
2 parts_err
2 notify_sw
2 write sw
2 pad
2 ht_thread
2 padi
2 pad
2 used_thread 3 fp
3 bp
t fixed bin (36) uns
bit (36) aligned, bit (36) aligned bit (1) aligned.
fixed bin, fixed bin, bit (36) aligned bit (18).
fixed bin.
(vtoc_buffer_seg.n_hash_buckets) bit (18) aligned,
(vtoc_buffer_seg.n_bufs) aligned like vtoc_buf_desc,
(vtoc_buffer_seg.n_bufs) aligned like vtoce_buffer:
/* Owner */

* For lock */
/* $\mathrm{DN}=$ = notify on unlock */
/* Number of full VTOCE buffers */
* Number of hash table buckets */
/* Mask for hash algorithm */
/* Absolute address of vtoc_buffer_seg */
* Offset of buf_desc *
* Offset of buf */
/* Offset of hash table */
/* For now */
* Global lock for vtoc buffers */
** Constant to add to part index to form wait event */
aligned based (vtoc buf descp)
bit (18) unal. bit (18) unal, $\quad$ * Offset to PVTE
fixed bin (17) unal. $/ *$ VTOcE Index */
bit (18) unal.
bit (3) unal.
bit (3) unal.
bit (3) unal
bit (1) unal.
bit (1) unal
bit (7) unal.
bit (7) unal.
bit (18) unal.
bit (18) unal.
aligned,
bit (18) unal
bit (18) unal;
/* Offset to first part descriptor for this buffer */
/* Mask of parts used or os */
1* Mask of parts out-of-service */
/* Mask of parts with I/O errors (hot) */
/* ON => notify requied on $1 / 0$ completion */
/* ON => write $1 / 0$ */
/* Offset of next entry in hash table */
/* Used list thread */
/* Forward pointer */
/* Backward pointer */

1 vtoce_buffer 2 words

N PARTS PER VTOCE VTOCE BUFFER SIZE
aligned based.
(3 * 64) bit (36) aligned;
fixed bin int static options (constant) init (3); fixed bin int static options (constant) init ( $3: 64$ );

